

WAECO Intro to AECs talk part 1

Based largely on material learned in Rami Grossberg's Model Theory II course.

1st order logic: All families with theory T in language L ,
study $(\text{Mod}(T), \leq)$
 $\uparrow$ elementary substructure.

Call these elementary classes.

Lots of useful tools:

- compactness
- LS theorems ($\uparrow$ & $\downarrow$)
- Chain theorems ($\langle M_i : i < \delta \rangle$ cts & inc,
then $\bigcup M_i \in \text{Mod}(T)$,
 $M_j \preceq \bigcup M_i$, and
 $M_j \preceq N \forall j \Rightarrow \bigcup M_i \preceq N$)

But there are other logics & classes of interest, e.g.

⊛ Infinitary logic: Models of some $\mathcal{Q} \in L_{\infty, \omega}$ perhaps

$(\text{Mod}(\mathcal{Q}), \leq_F)$ where $\mathcal{Q} \in L_{\kappa^+, \omega}$

(i.e. $\mathcal{Q}$ allows $\wedge$ and $\vee$ of size $< \kappa^+$), $|L| \leq \lambda$,
and F a fragment of $L_{\kappa^+, \omega}$ containing $\mathcal{Q}$ with $|F| \leq \lambda$.

(i.e. $F \subseteq \text{Form}(L_{\kappa^+, \omega})$ where

- F closed under f.o. connectives $\wedge, \vee, \neg, \forall x, \exists y$

- If $\mathcal{Q}(x) \in F$, $\tau(y) = \langle \tau_1(y), \dots, \tau_n(y) \rangle$
where τ_i terms, then $\mathcal{Q}(\tau(y)) \in F$

- F is closed under subformulas.

so ' F is well enough closed'

$$\textcircled{\star} \quad EC(T, \Gamma) = (\{M \models T : M \text{ models all } \varphi \in \Gamma\}, \leq).$$

In these, compactness fails - so if we move to a more general framework that also covers these examples, we need to weaken our assumptions.

Definition: An abstract class (AC) is a pair $\underline{K} = (K, \leq_K)$ where

- K is a class of models in some language $L(K)$
- $\leq_K$ is a partial ordering on K refining submodel, i.e.
 $M \leq_K N \Rightarrow M \subseteq N$

- 1st iso axiom: K is closed under isomorphism: i.e.
 if $M \in K$, $N \cong M$, then $N \in K$.

- 2nd iso axiom: $\leq_K$ respects isomorphism, i.e.
 given $M_1 \leq M_2$ and $M_2 \stackrel{f}{\cong} N$
 then $f[M_1] \in K$ and $f[M_1] \leq N$

$$\begin{array}{ccc} M_2 & \xrightarrow{\cong f} & N \\ \downarrow \text{V/K} & & \downarrow \text{V/K} \\ M_1 & \xrightarrow{f} & f[M_1] \end{array}$$

These assumptions are still too weak to capture much of the theory of the last two examples. So we have...

Definition: An abstract elementary class (AEC) is an AC $\underline{K} = (K, \leq_K)$ that also satisfies

- Cohesiveness: If $M_1, M_2, M_3 \in K$ and $M_1 \leq_K M_3$,
 $M_2 \leq_K M_3$, and $M_1 \subseteq M_2$, then $M_1 \leq_K M_2$.

— TV chain axioms: If $\langle M_i : i < \delta \rangle$ a $\leq_K$ -inc cts $\leq_K$ -chain of models in K where δ is limit, then

- $\bigcup_{i < \delta} M_i \in K$
- $M_j \leq_K \bigcup_{i < \delta} M_i$ for all $j < \delta$
- If $M_j \leq_K N$ for all j , then $\bigcup_{i < \delta} M_i \leq_K N$.

— LS Axiom: There is an infinite cardinal $LS(K)$ such that
 $\forall M \in K \forall A \subseteq |M| \exists N \in K$ s.t. $N \leq_K M$, $A \subseteq |N|$,
 and $||N|| \leq |A| + LS(K)$.

We denote the least such cardinal by $LS(K)$.

— (sometimes) $\forall M \in K \quad ||M|| \geq LS(K)$.

Examples: • $(\text{Mod}(T), \leq)$, $LS(K) = |L(T)| + \aleph_0$

• $(\text{Mod}(\varphi), \leq_F)$ where $\varphi \in L_{\lambda^+, \omega}$, $|L|, |F| \leq \lambda$
 then $LS(K) = \lambda$

• $EC(T, T)$, $LS(K) = |L(T)| + \aleph_0$

• Lots of specific algebraic examples:

— groups with $\leq_K = \subseteq$

or $\leq_K = \text{'pure subgroup'}$

$$\left(i.e. H \leq_K G \iff \forall n \in \omega \forall a \in H \left[G \models \exists x (a = x^n) \Rightarrow H \models \exists x (a = x^n) \right] \right)$$

— locally finite groups

— nil rings $(\forall n \in \mathbb{N} \exists n \in \omega \ x^n = 0)$

} with $\leq_K = \subseteq$, $LS(K) = \aleph_0$

— modules over a ring with $\leq_K = \subseteq$, $LS(K) = |R| + \aleph_0$

— Well orders of type $\leq \lambda^+$ with $\leq_K = \text{'initial segment'}$, $LS(K) = \lambda$

— models over a ring with $\leq_K = \text{'pure submodule'}$, $LS(K) = |R| + \aleph_0$.

Non-examples:

- well orders with $\leq_K = \subseteq$ (chain breaks )
- well orders with $\leq_K$ initial segment (LS axiom breaks) not well founded

Notation: • Given λ cardinal, $K_\lambda = \{M \in K : \|M\| = \lambda\}$

• $I(K, \lambda) = |K_\lambda / \cong|$ ($\leq 2^{LS(K) + \aleph_0 + \lambda}$)

Question: What can $I(K, \lambda)$ look like?

Note $\{\lambda \geq LS(K) : I(K, \lambda) \neq 0\} = [LS(K), *]$ for some $*$ or $*$ = ∞
by the LS axiom.

$*$ $< \infty$ is possible, but

Fact (Shelah): If $*$ $\geq \beth_{(2^{LS(K)})^+}$, then $*$ = ∞ .

Fact: There are examples of $*$ = $\beth_\alpha$ for all $\alpha < (2^{LS(K)})^+$.

Considering Morley's / Shelah's categoricity theorems,

Question: what can $\{\lambda : I(K, \lambda) = 1\}$ look like?

Big open problem:

Shelah's categoricity conjecture: If $\mu_0 = \beth_{(2^{LS(K)})^+}$ and $\exists \lambda_0 \geq \mu_0$

such that K is categorical in λ_0 , then K is categorical in all $\lambda \geq \mu_0$.

or maybe use a different bound for μ_0 in terms of $LS(K)$.

Lifting basic notions to an AEC setting:

Rather than elementary embeddings ...

Definition: Given $M, N \in K$, $f: M \rightarrow N$ is a K -embedding iff
 $\exists N^* \leq_K N$ such that $f: M \cong N^*$ (i.e. $f: M \cong f[M] \leq_K N$)

Rather than types ...

Definition: K an AEC.

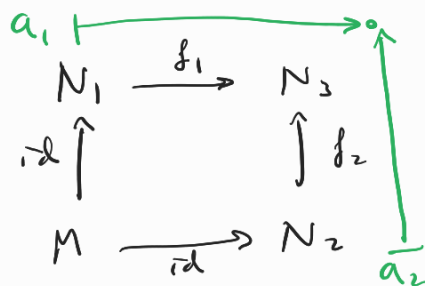
- $K^3 = \{(M, N, a) : M, N \in K, a \in N, M \leq N\}$.

- Given $(M_\ell, N_\ell, a_\ell) \in K^3$ for $\ell = 1, 2$

$(M_1, N_1, a_1) \sim (M_2, N_2, a_2)$ iff $M_1 = M_2$ and

$\exists N_3 \in K \exists f_\ell: N_\ell \xrightarrow{M_1} N_3$ K -embeddings such that

$$f_1(a_1) = f_2(a_2)$$



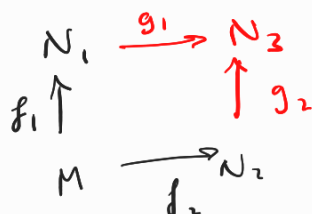
- $\approx$ is the transitive closure of $\sim$

- Given $(M, N, a) \in K^3$, $\text{gtp}(a/M, N) = [(M, N, a)]_{\approx}$.

In many cases, $\sim = \approx$.

Definition: Suppose K is an AC. We say K has the amalgamation property if whenever we have $f_\ell: M \rightarrow N_\ell$ for $\ell = 1, 2$,

$\exists N_3 \exists g_\ell: N_\ell \rightarrow N_3$ such that $g_1 \circ f_1 = g_2 \circ f_2$

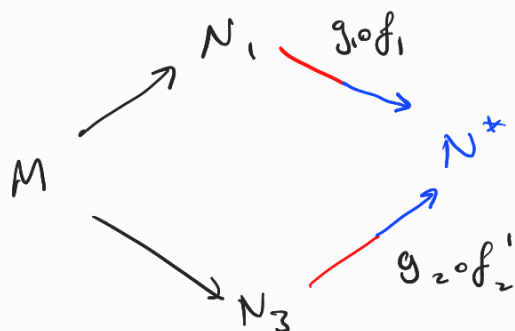
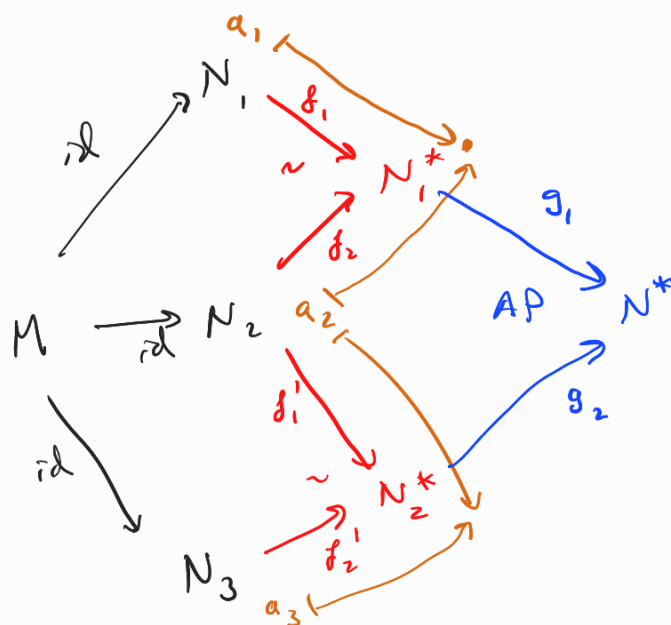


Proposition: If K has AP, then $\sim = \simeq$.

proof: Suppose $(M_1, N_1, a_1) \sim (M_2, N_2, a_2) \sim (M_3, N_3, a_3)$.

Let $M = M_1 = M_2 = M_3$.

Sketch:



$$g_1 \circ f_1(a_1) \stackrel{\sim \text{def}}{=} g_1 \circ f_2(a_2) \stackrel{AP}{=} g_2 \circ f_1'(a_2) \stackrel{\sim \text{def}}{=} g_2 \circ f_2'(a_3)$$

Definition: Let K be an AEC, $M \in K$.

$$S(M) = \{g \upharpoonright N(a/M, N) : M \leq N, a \in |N|\}$$

Fact: $|S(M)| \leq 2^{\|M\|}$

(wlog $\|N\| = \|M\|$ by LS Ax & def of $\sim$,
 $\leq 2^{\|M\|}$ choices for $|N| \setminus |M|$,
 $\leq 2^{\|M\|}$ choices for interpreting $L(M)$ in N ,
 $\leq \|M\|$ choices for A .)

Definition: If $p = \text{gfp}(a/M, N)$, we say p is realized in N .

Definition: M is λ -saturated if $\forall M_0 \leq M$ s.t. $\|M_0\| < \lambda$,
 $\forall p \in S(M_0)$, p is realized in M .

M is saturated if $\|M\|$ -saturated.

These still behave as we expect.

Proposition: Suppose K is an AEC, $\lambda > LS(K)$, $K_{<\lambda}$ has AP,
and $M \in K_\lambda$. Then

M is saturated $\iff \forall N \leq M \forall N' \geq_K N$ s.t. $\|N'\| < \|M\|$

$$\exists f: N' \xrightarrow[N]{} M$$

M is model homogeneous.

proof ideas: Like 1st order, but be careful.

sat $\iff$ m-h $\iff \forall \mu < \lambda$ M is m-h for $\|N\| = \|N'\| = \mu$

like 1st order,
but there is no notion of (M, N) -map in AECs,
(all sets are models), so we have all maps between
'larger' models.

Other important notions:

Definition: We say an AC K has the joint embedding property (JEP) if
 $\forall M_1, M_2 \in K \exists N \in K \exists f_1: M_1 \rightarrow N$.



Fact: $(\text{Mod}(T), \leq)$ has JEP $\Leftrightarrow T$ is complete.

($\Rightarrow$ if not complete, get $M_1 \neq M_2$)

$\Leftarrow$ use monster model, or large saturated model, as N).

Definition: We say an AC K has no maximal models (NMM) if

$$\forall M \in K \exists N \in K \quad M \subsetneq N.$$

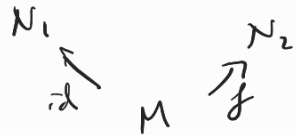
Fact: If K has AP, JEP, NMM, then K has monster models similar to 1st order logic, i.e. $\forall \mu > \text{LS}(K) \exists M^* \in K$ s.t.

- All models of card $< \mu$ embed into M^*
- M^* is μ -m.h. (i.e. m.h. holds for $\|N\| = \|N'\| < \mu$)
- $\forall N, N' \leq_K M^*$ s.t. $\|N\| = \|N'\| < \mu$, $\forall f: N \cong N'$
 f extends to an automorphism of M^* .

Note: Existence of such M^* recovers AP, JEP, NMM

proof: JEP, NMM easy

AP:



$M, N_1 \leq M^*$ wlog
 get $g: N_2' \cong N_2$ extending f . By μ -m.h.,

$$\exists h: N_2' \xrightarrow{M} M^*$$

$$\begin{array}{ccc} N_1 & \xrightarrow{id} & M^* \\ id \uparrow & & \uparrow h \circ g^{-1} \\ M & \xrightarrow{f} & N_2 \end{array} //$$

Note: - With a monster model, types $gtp(a/M, N)$ over small sets become orbits under automorphisms fixing M

- The notion of types agrees with the first order notion: if $a \in N_2 \geq_K^M e$,
 $tp(a/M_1, N_1) = tp(b/M_1, N_1) \Leftrightarrow gtp(a_1/M_1, N_1) = gtp(a_2/M_2, N_2).$

END OF PART I